On the Usefulness of Liveness for Garbage Collection and Leak Detection

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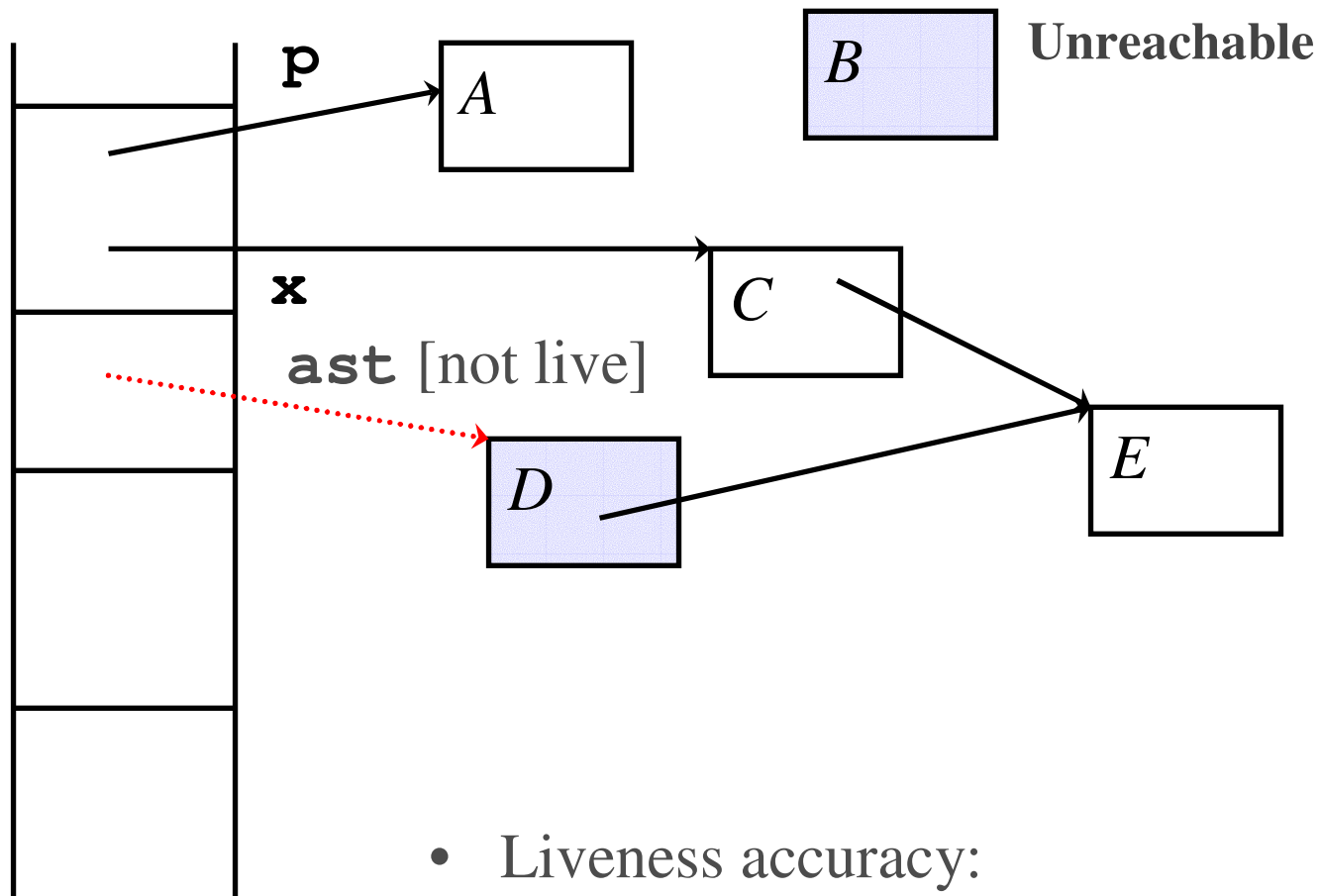
ECOOP June 2001 Budapest, Hungary

What is Liveness?

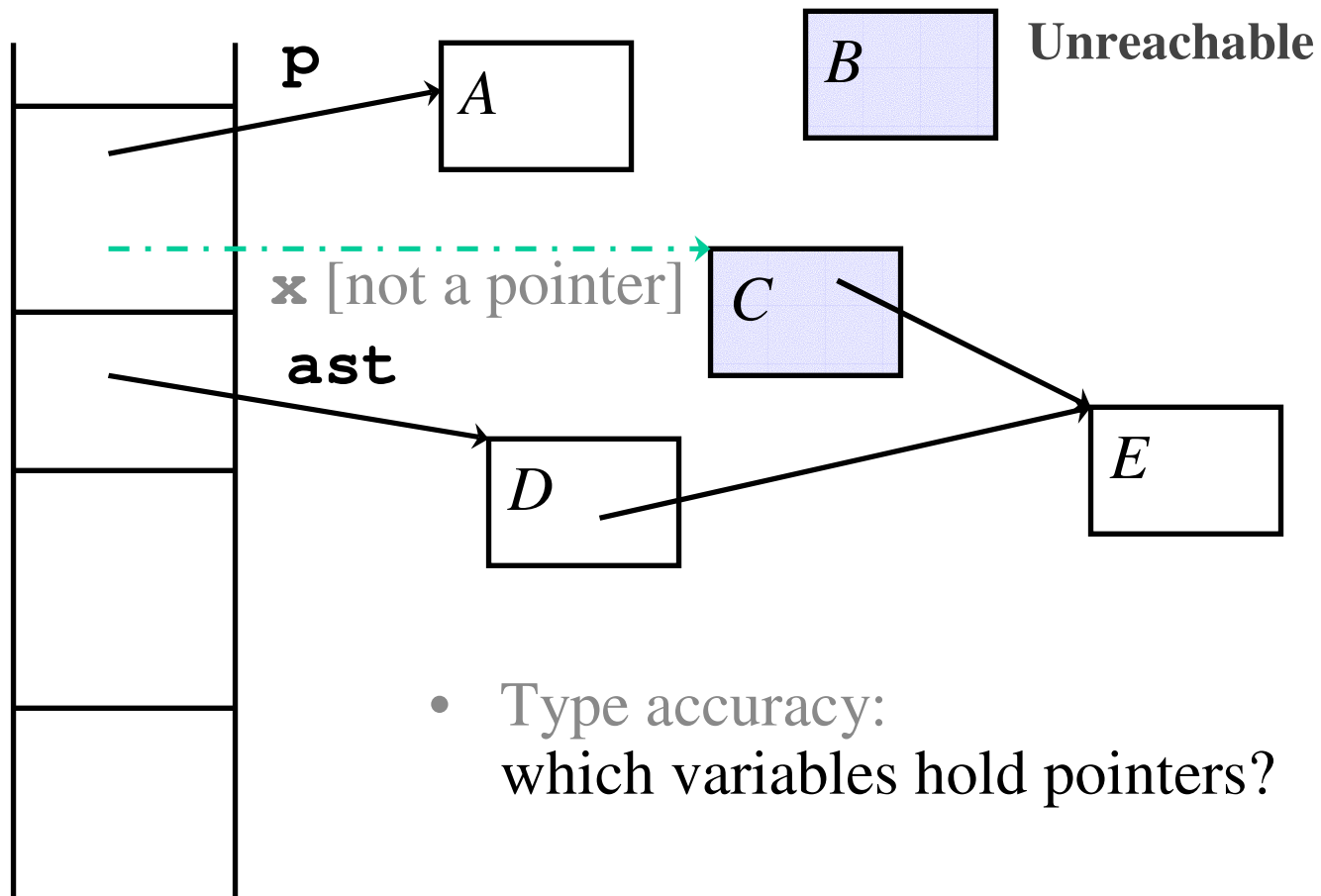• A variable is *live* if its value will be used in the future.

```
…                                      ast

Tree *ast = parse();

Cfg *cfg = translate(ast);

⟨  code that does not use the value
    of ast  ⟩

...
```

Accuracy and Reachability

Accuracy and Reachability

p

A

B   **Unreachable**

x

C

ast [not live]

D

E

- Liveness accuracy:
  which variables are live?

Accuracy and Reachability

p

A

B   **Unreachable**

x [not a pointer]   C

ast

D

E

- Type accuracy:
  which variables hold pointers?

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Accuracy and Reachability

p

A

B   **Unreachable**

x [not a pointer]   C

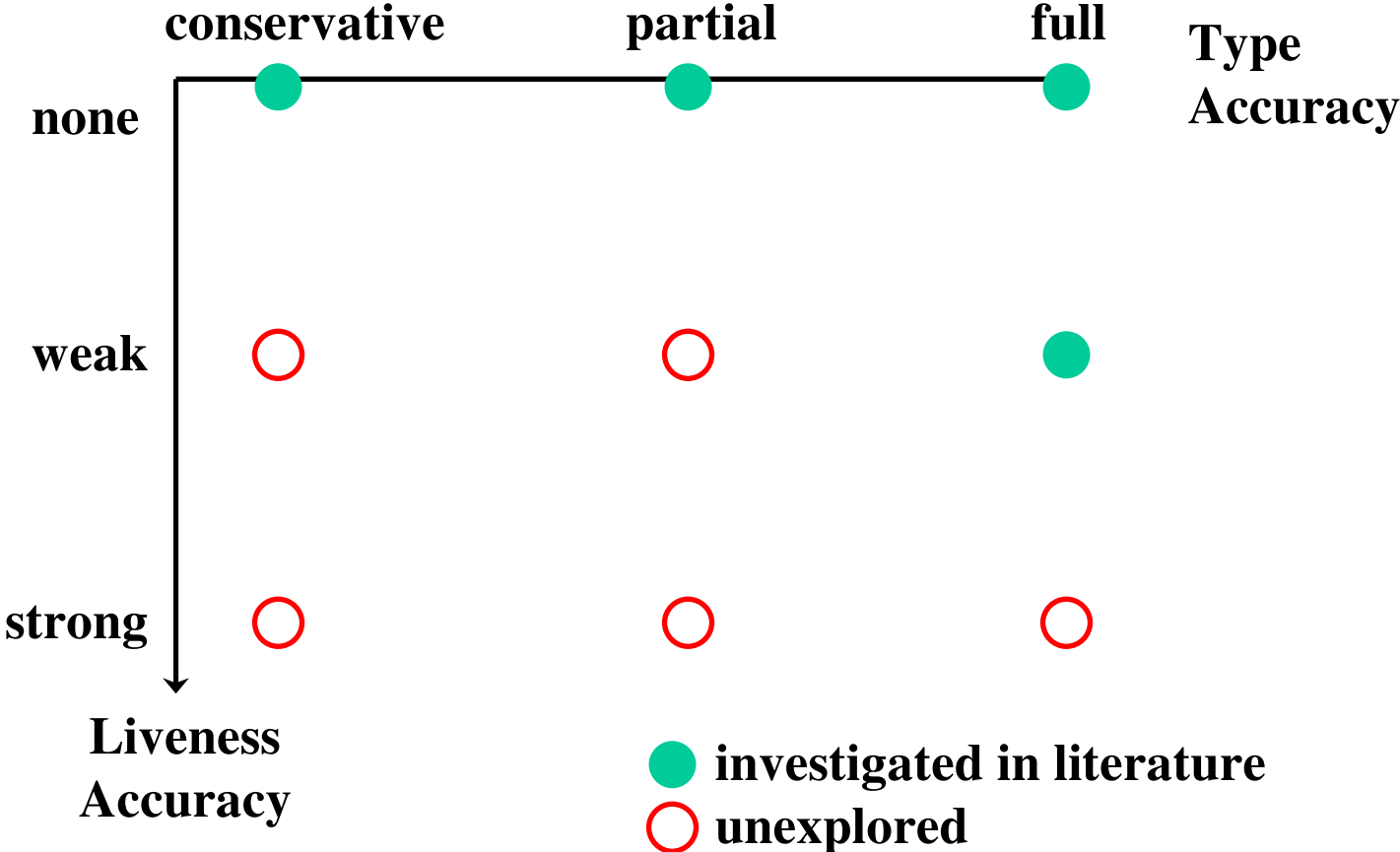**ast** [not live]

D   E

- Type accuracy:
  which variables hold pointers?
- Liveness accuracy:
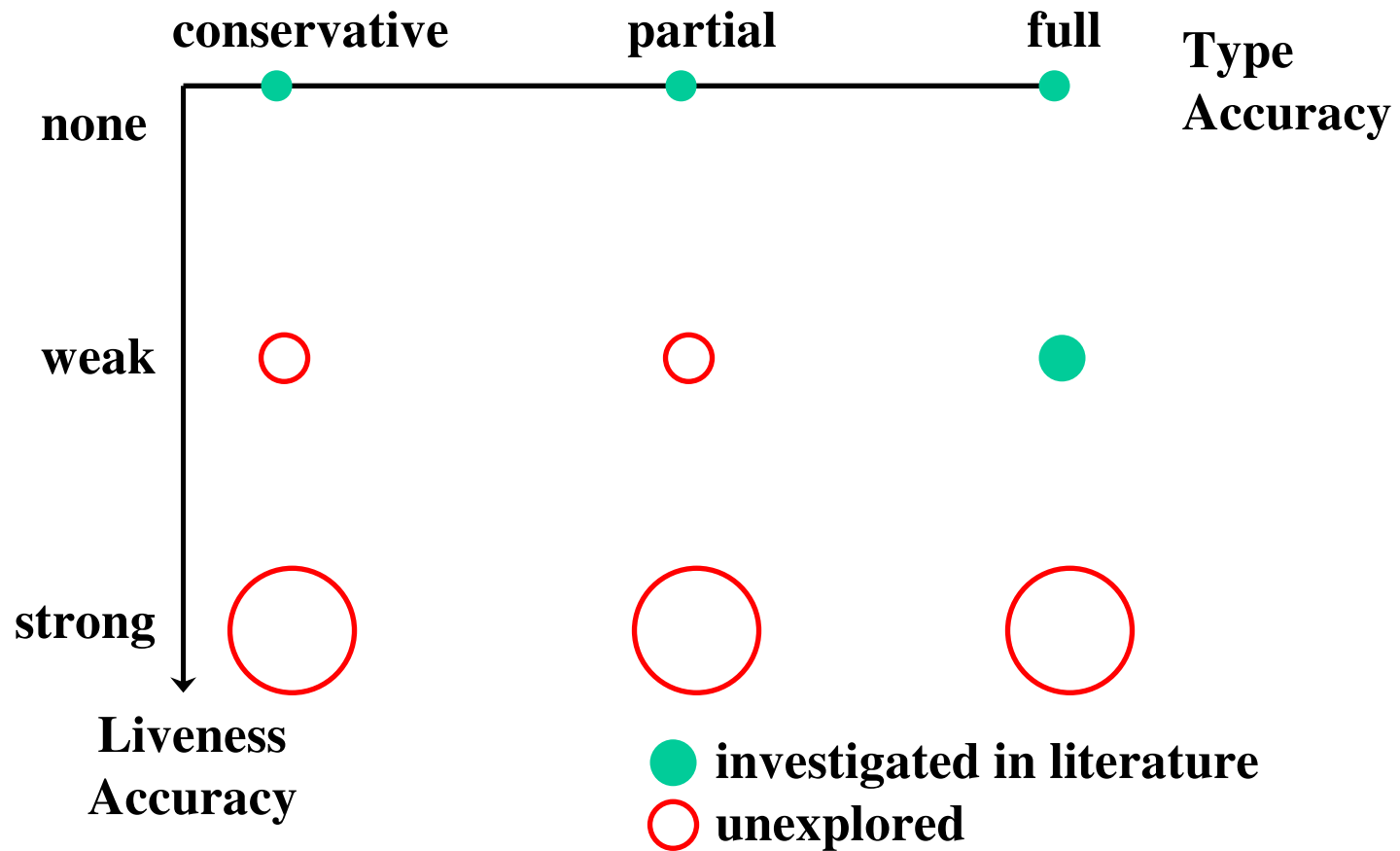 which variables are live?

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The Questions

- Can liveness accuracy benefit reachability traversals?
  - Useful for garbage collection?
  - Useful for leak detection?
- What kind of analysis is necessary?

Motivation

Results in a Nutshell

Outline

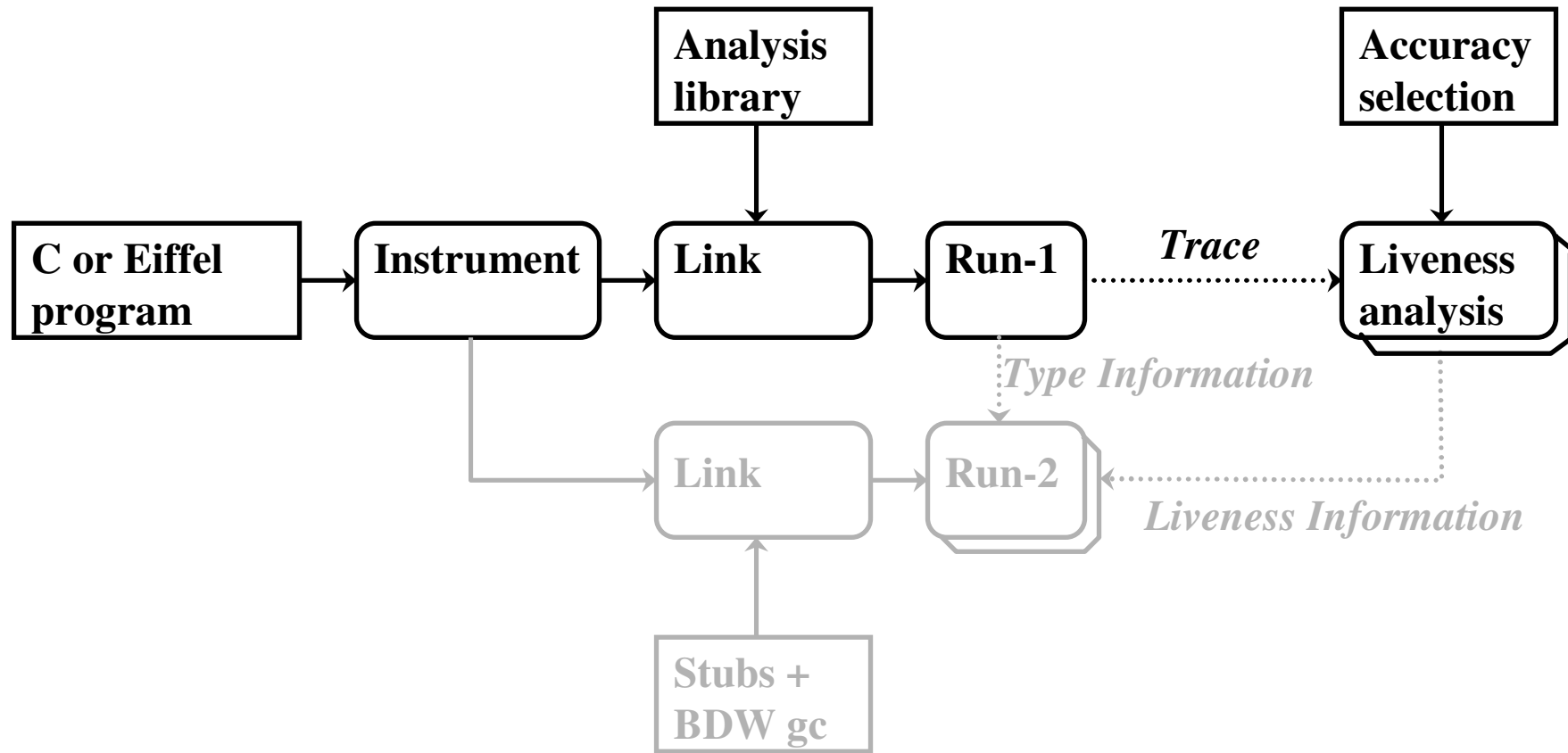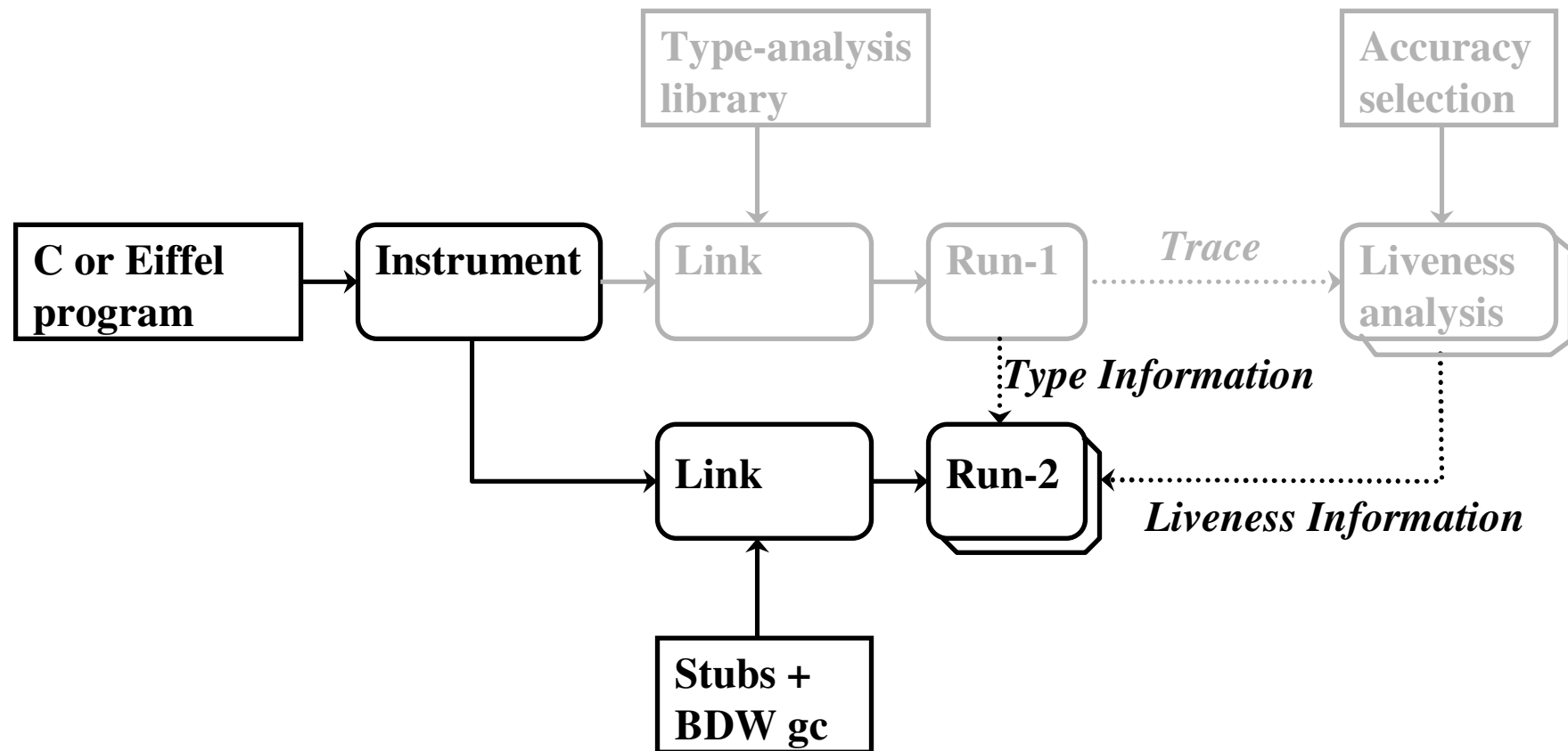| Introduction	• Motivation
	• Preview of results
Methodology	• Obtaining liveness information
	• Metrics
Results	• Reachable heap given various accuracy schemes
Conclusion	• Related work
	• Summary

Liveness Approaches

- Static analysis
  - Compiler-analysis of source code
  - Disadvantage: difficult, cost + benefits unclear

- Dynamic analysis
  - Trace-based analysis
  - Disadvantage: two runs needed, limit study

Infrastructure for Experiments

Infrastructure for Experiments

Generating the Trace

Original Code	Instrumented Code	Trace

```x = malloc(4*N);```
 	
```i = 0;```
 	
```while(i < N){```
  ```y = x + 4*i;```
 	
  ```*y = i;```
 	
  ```i++;```
 	
```}```
	```x = malloc(4*N);```
	```note_allocation();```
	```note_assign(&x);```
	```i = 0;```
	```note_assign(&i);```
	```while(i < N){```
	  ```y = x + 4*i;```
	  ```note_assign(&y,&x,&i);```
	  ```*y = i;```
	  ```note_use(&y);```
	  ```note_assign(y,&i);```
	  ```i++;```
	  ```note_assign(&i,&i);```
	```}```
		…
		allocation(91)
		assign(x)
		assign(i)
		assign(y, x, i)
		use(y)
		assign(91.0, i)
		assign(i, i)
		assign(y, x, i)
		use(y)
		assign(91.4, i)
		…

Analyzing the Trace

Trace	Simulated Liveness State	Resulting Information
	x y i	

…
allocation(91)
assign(x)
assign(i)   T1: {*x, i*}
assign(y, x, i)
use(y)
assign(91.0, i)
assign(i, i)   T1: {*x, i*}
assign(y, x, i)
use(y)
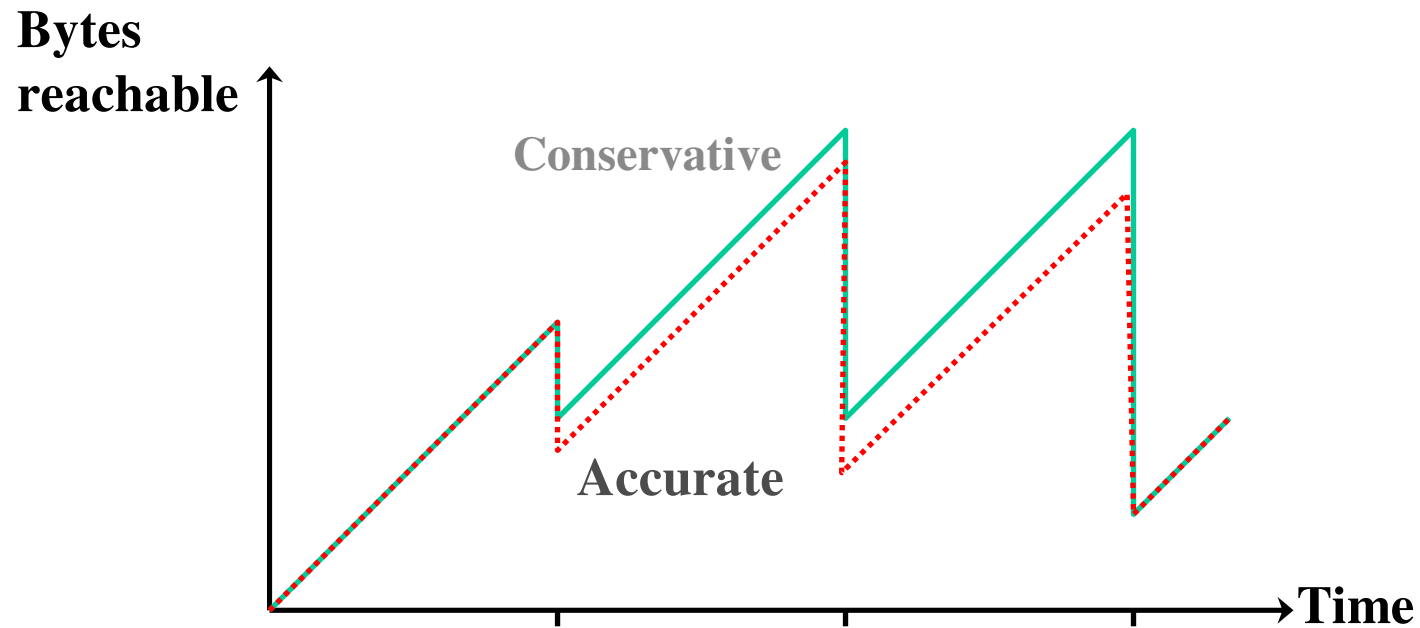*assign(91.4, i)*   T2: {}
…

"A variable is *live* **if its value will be used in the future."**

Usefulness Metric for Accuracy

Bytes reachable

Conservative

Accurate

Time

Reachability traversal: 1   2   3
Reachability reduction: 10%   20%   0%

⇒ **The accurate scheme reduced reachability by 10% on average.**

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Benchmarks

Name	Language	Lines of Code	Total allocation [Bytes]	Author/Source
Programs written with GC in mind:				
gctest3	C	85	2 200 004	Bartlett
gctest	C	196	1 123 180	Bartlett
bshift	Eiffel	350	28 700	Hirzel
erbt	Eiffel	927	222 300	Durian
ebignum	Eiffel	3 137	109 548	Hillion
li	C	7 597	9 030 872	Spec95
gegrep	Eiffel	17 185	106 392	Bezault
Programs with explicit deallocation:				
anagram	C	647	259 512	Austin
ks	C	782	7 920	Austin
ft	C	2 156	166 832	Austin
yacr2	C	3 979	41 380	Austin
bc	C	7 308	12 382 400	Austin
gzip	C	8 163	14 180	GNU
ijpeg	C	31 211	148 664	Spec95

Usefulness of Liveness

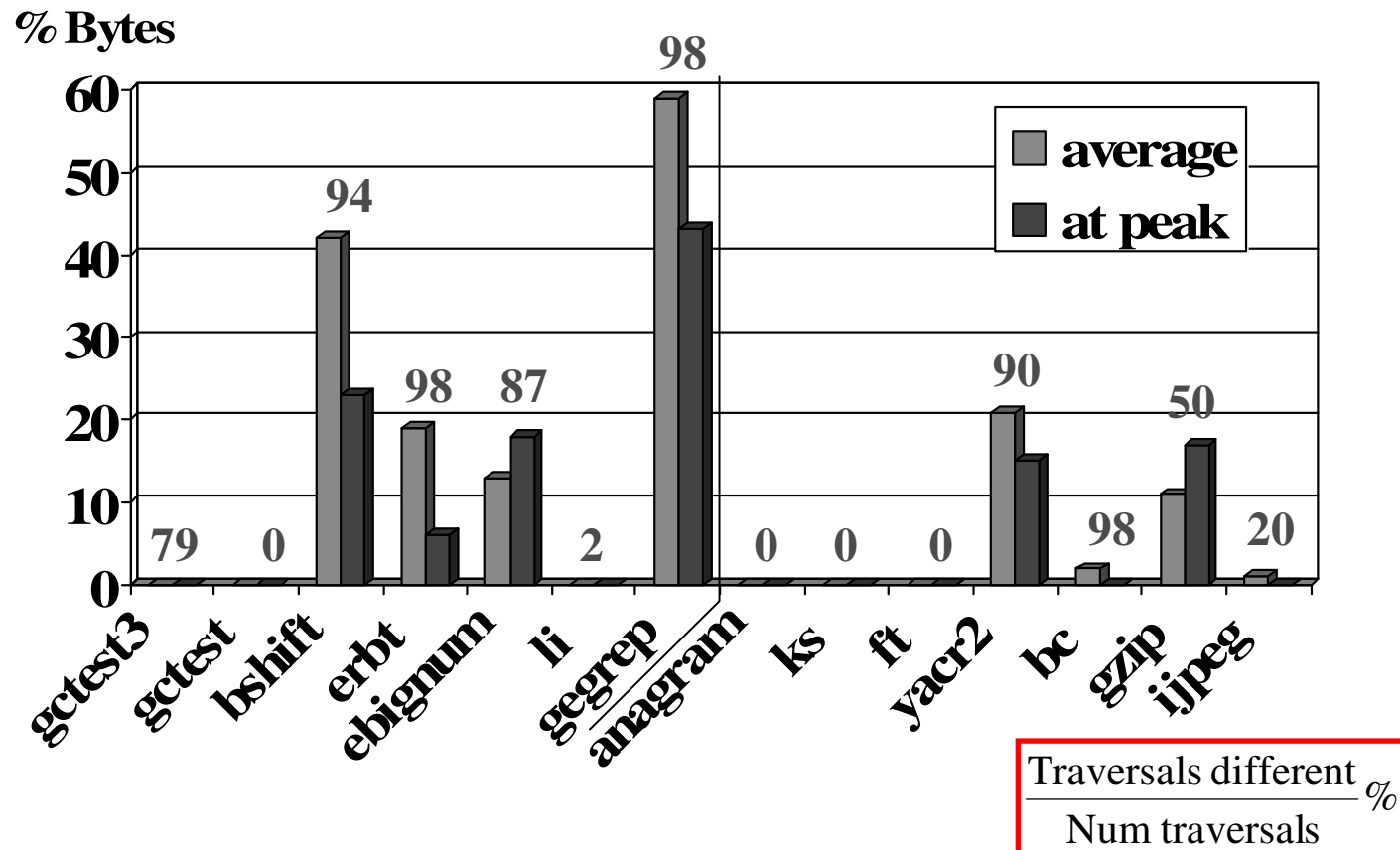**Reachability reduction with strongest liveness**

Usefulness of Liveness

Reachability reduction with strongest liveness

% Bytes

Legend: average, at peak

Data labels: 79, 0, 94, 98, 87, 2, 98, 0, 0, 0, 90, 98, 50, 20

X-axis: gctest3, gctest, bshift, erbt, ebignum, li, gegrep, anagram, ks, ft, yacr2, bc, gzip, ijpeg

$$\frac{\text{Traversals different}}{\text{Num traversals}}\%$$

20

Different Levels of Liveness

% Bytes

Reachability reduction

Legend:
- stack intra scalars
- stack inter scalars
- stack inter all
- s.+glob intra scalars
- s.+glob inter scalars
- s.+glob inter all

Y-axis: 0, 10, 20, 30, 40, 50, 60

X-axis: gctest3, bshift, erbt, ebignum, li, gegrep, yacr2, bc, gzip, ijpeg
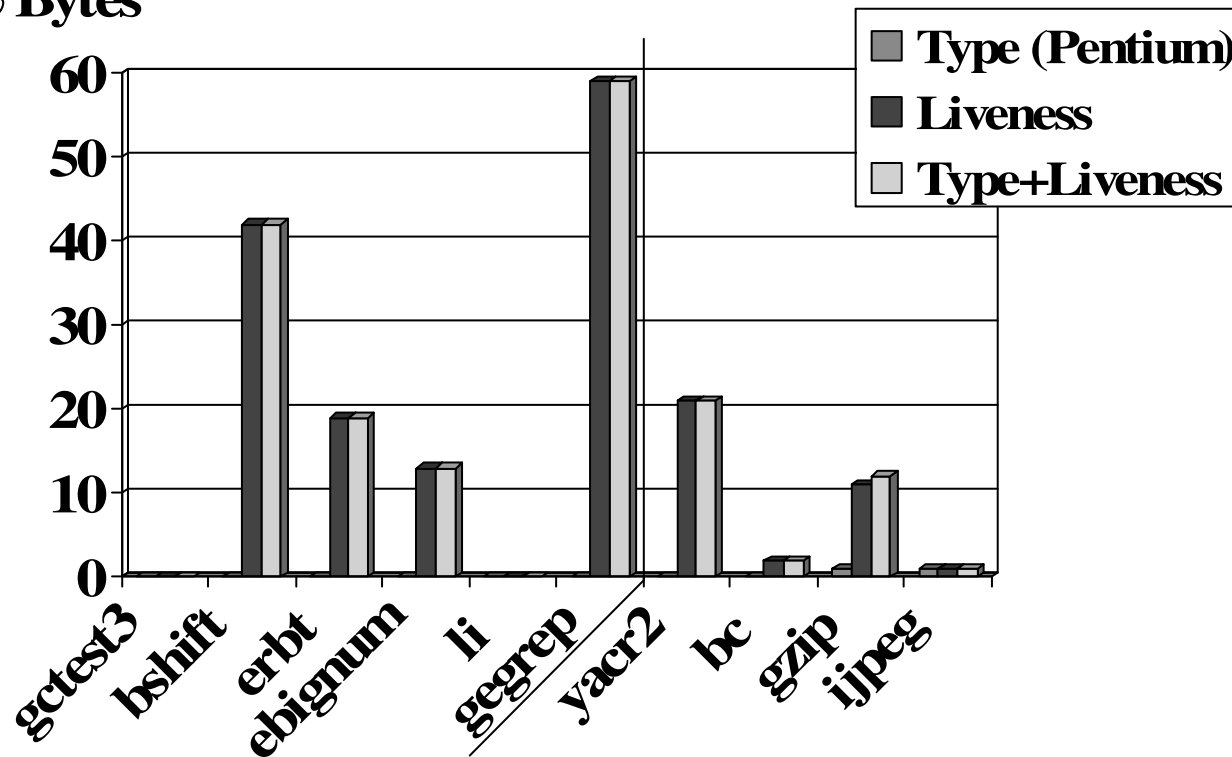
Type versus Liveness Accuracy

Reachability reduction

Validation

- Comparing liveness information found in different runs
  - For how many locations did the obtained liveness information differ?

Benchmark	Stack % different	Global % different
gegrep	0.7	0.0
yacr2	2.7	0.0
gzip	1.3	2.2

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Related Work

- Evaluating Accuracy
  - Hirzel, Diwan: On the type accuracy of garbage collection. ISMM 2000.
  - Shaham, Kolodner, Sagiv: On the effectiveness of GC in Java. ISMM 2000.
- Implementing Accuracy
  - [Bartlett1988] [DiwanMossHudson1992] [SmithMorrisett1998] [Zorn1993] [AgesenDetlefsMoss1998] …

Summary

- Liveness accuracy can be very useful for reachability traversals.

- Strong analyses are necessary to reach significantly fewer Bytes.

- Type accuracy was not very useful in these experiments.