

Robust Scripting via Patterns

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Setting

- Thorn language
 - IBM and Purdue project, now in stasis
- Dynamic Languages
 - No static types
- Concrete Data Structures
 - Lists, records, objects / datatypes
- Imperative languages
 - But emphasis on declarative/functional

Related Work

- SNOBOL4 (1966)
- ML, ISWIM, Hope, Haskell, F#, Scala, Kotlin
- Scheme, Newspeak, Python, Converge, OMeta
- OCaml, JMatch
- Views, Tom, Matchete

Plan

- Pattern Language
 - Some fancy patterns
 - First-class Patterns
- Integration with Thorn
 - Patterns used everywhere
 - Some interactions with standard control flow
- Usage
 - **Do** Thorn programmers do what they **can** do?

Patterns (in the ML Sense)

- Match a **subject** value against a **pattern**
 - Can **FAIL**
 - Can **SUCCEED** and bind some **variables**

Name	Subject	Pattern	Result	Bindings
Variable	[1,2,3]	x	succeed	x=[1,2,3]
List	[1,2,3]	[x]	fail	
Wildcard	[1,2,3]	[x,_,_]	succeed	x=1
Head/Tail	[1,2,3]	[x, y...]	succeed	x=1, y=[2,3]
Literal	[1,1]	[x, 1]	succeed	x=1
Value	[1,1]	[x,\$x]	succeed	x=1
Record	<a=1,b=2,c=3>	<a=x, b>	succeed	x=1, b=2

How Much Are They Used?

- Corpus:
 - 24K lines of code
 - Most of the Thorn code in existence
- Coders
 - Bard (60%), skilled (30%), novices (10%)
- Purposes
 - Some examples of Good Thorn Style
 - Some one-shot programs to throw away
- **This Is Not Science**
 - Literary Analysis, maybe
- Negative results may be interesting too

Part I: Control and Patterns

Control Structures and Patterns

- **Design Principle:** Put patterns wherever they might make sense
- **Design Principle:** Patterns should be allowed wherever variables are bound to arbitrary values
 - If it makes sense
 - Deal with failure somehow
 - *E.g.* Formal parameters can be patterns

Binding Statement

- Binding statement (LISP/ML let):
 - `x = [1 , 2 , 3]`
- With pattern, it's destructuring
 - `[a , b , c] = [1 , 2 , 3]`
 - Exception if fails
- Usage: 3% of bindings have interesting pattern
 - Bard prefers defensive programming

Scopes

- **Design principle:** pattern matches introduce variables into the scope that will be executed iff the match succeeds.
- Match Operation: $E \sim P$
 - returns true on success, false on failure
 - Produces bindings in right scope
- But what's the right scope?
 - Depends on context...

if statement

- `if (L ~ [x])
 use(x);
else
 xUndefined();`
- We support
 `if (A ~ P && B ~ Q && C ~ R)`
 – (But not general propositional logic)
- 37% of if's have matches
- *(There's a match statement too, but much less used than 'if')*

Patterns and while

- while: bindings in test can be used in body

```
while(R ~ <x>)  
    R := munge(R, x);  
xUndefined();
```

Patterns and until

- Until: bindings in test can be used **after** body
 - `until(x.spouse ~ (!null && y))`
 `x.date();`
 `fileJointly(x,y)`
 - Precisely expresses "look for something"
- Rarely used (<1%)
 - Searching comprehensions and recursion are favored.
 - Thorn bias: Most whiles were `while(true)` in actor bodies

Patterns and Control, reprise

- There's value to making patterns aware of control:
 - if, for: 40%
 - fun, lambda: 20%
 - let, while, until: 1-3%

Part II: Fancy Patterns

Kinds of Patterns

- Common Patterns
 - Most patternly languages have these
 - wildcard, variable, literal, list, ...
 - 82% of Thorn patterns are common
 - Count of syntax tree nodes
 - Not counting variables
- Fancy Patterns
 - Few languages have any of these
 - Fewer have all of them.
 - 18% of Thorn patterns are fancy
 - Let's see a couple...

Fancy Pattern: Type Test

- **General form:** $P : T$
 - matches a value of type T
 - which must also match pattern P
 - And binds what P does
- **Idiom:**
 - `fun f(x:int) = x+3;`
- **Usage:** 3.5% of all patterns

Fancy Pattern: Boolean Combinations

Pattern	Matches	Binds	Usage
P && Q	if both P and Q match	Everything bound by P or Q (disjoint)	3%
P Q	if either P or Q matches	Everything bound by both P and Q	0.2%
!P	if P fails	nothing	0.1%

&& is useful

- Pattern: `x && [y, z...]`
 - Matches a nonempty list
 - Binds the whole list (x), the head (y) and tail (z)
- *as* construct in pattern-bearing languages
 - "Get a whole value and its parts"
- Trans-*as* usage:
 - `[_..., 1, _...] && [_..., 2, _...]`
 - Matches a list containing 1 and 2 in either order
- About 3% of patterns involve &&
 - Mostly for the *as* idiom.
 - No popular idioms for `| |` and `!`
 - A good idiom makes a pattern operator popular.

Internal Matches

- General Form: $E \sim P$
 - Succeeds if value of E matches P
 - Binds what P does
 - Can appear inside of patterns
 - Usage: 3.5%
- Example: $[x] \ \&\& \ f(x) \ \sim \ [y, z]$
- Swiss Army Construct
 - E.g. optional field foo, defaulting to 22:
 $\langle \text{foo}=\text{x} \rangle \ \|\ \|\ 22 \sim \text{x}$

Part III: First-Class Patterns

- Fanciest of all the fancy patterns.

First-Class Patterns

- First-class **functions** are amazingly useful
 - One of the top N ideas in programming languages

First-Class Patterns

- First-class **functions** are amazingly useful
 - One of the top N ideas in programming languages
- First-class **patterns** are a bit cool
 - One of the top N^3 ideas in programming languages

Why abstract patterns?

- Summing binary trees
- Object/datatype representation:

```
fun sum(Fork(l, x, r)) = sum(l) + x + sum(r);  
  | sum(Leaf(x))      = x;
```

(This is the nicest code in the universe)

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- List representation:

```
fun sum([l, x, r])    = sum(l) + x + sum(r);  
  | sum([x])          = x;
```

(this is also the nicest code in the universe)

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```

- Are we not computer scientists?
 - And do we not abstract reflexively?

Pattern Expression, part 1

- Pattern Abstraction:
 - A **value** (not a **pattern**).
 - `pat [x,y] = [x,$x,y]`
 - `x,y` are **outputs** not **inputs**.
 - `x,y` are scoped inside the expression
- Pattern Application
 - `E[r,s]` is a **pattern**
 - `r,s` are **subpatterns**
 - Appears in pattern context: `somelist ~ E[r,s]`

```
E = pat[x,y] = [x,$x,y]
```

```
L = [3, 3, 4]
```

```
if (L ~ E[a,b]) assert(a==3, b==4)
```

```
if (L ~ E[a,9]) fails()
```

Sum with Representation Parameter

- Representation pattern

```
rp = <fork=fpat, leaf=lpattern>
```

```
– rp.fork[l,x,r] matches a fork node
```

```
– rp.leaf[x] matches a leaf node
```

- Sum with explicit rp:

```
fun sum2(rp, rp.fork[l,x,r])
```

```
    = sum2(rp,l) + x + sum2(rp,r)
```

```
  | sum2(rp, rp.leaf[x]) = x
```

No longer the most beautiful code in the universe

Computing the Representation

```
// Guess representation of a tree...
```

```
fun rep([_,_,_] || [_]) = repList;  
  | rep(["Fork" || "Leaf", _...]) = repTaggedList;  
  | rep(x:Tree) = repTree;  
  | rep(<left,item,right> || <leaf>) = repRecord;
```

```
// Use it!
```

```
fun sum(rep(it).fork[l,x,r]) = sum(l) + x + sum(r);  
  | sum(rep(it).leaf[x]) = x;
```

Pattern Abstractions, parts 2-N

- More variations
 - pattern/constructor duality
 - inputs and outputs
- Late addition to language
 - We didn't get to use them much
- Nice new toy!

Conclusion

- There's a lot more to patterns than ML-style
 - $P \& \& Q$, $E \sim P$, $\text{pat}[x]=P$
- Patterns can be meshed with statements
 - `if(L~[x,y]) use(x,y);`
- If you have them, they will be used
 - happily!